

# ALGORITHMS FOR THE SOLUTION OF THE GENERALIZED EIGENVALUE PROBLEM

Z. BOHTE, J. GRAD

UDK: 521.643.5

INSTITUTE OF MATHEMATICS, PHYSICS AND MECHANICS  
JADRANSKA 19, LJUBLJANA, YUGOSLAVIA

In this paper we deal with generalized eigenvalue problem of matrix  $A(z)$  the elements of which are polynomials in  $z$ . The well known iterative methods of Muller, Newton and Laguerre for finding the zeros of function  $f(z) = \det A(z)$  are analysed. Decompositions of matrix  $A(z)$  and its derivatives are introduced in order to simplify the computations of the values of  $f(z)$  and its first and second derivatives. Comparative analysis gives some indicators about the rate of convergence, computer time and accuracy of the computed eigenvalues for each of the method used. The analyzed methods proved generally to be stable, economical and easily applicable. FORTRAN subroutines and an example of main calling programme are added for Muller's and Laguerre's methods which proved to be more efficient than Newton's.

ALGORITMI ZA REŠEVANJE POSPLOŠENEGA PROBLEMA LASTNIH VREDNOSTI. V članku je obravnavano numerično reševanje posplošenega problema lastnih vrednosti za matriko  $A(z)$  z elementi, ki so polinomi spremenljivke  $z$ . Primerjane so dobro znane iterativne metode Mullerja, Newtona in Laguerre za računanje ničel polinoma  $f(z) = \det A(z)$ . Vrednosti polinoma  $f(z)$  in njegovih odvodov so izračunane na osnovi razcepa matrike  $A(z)$ . Primerjava metod daje vpogled v hitrost konvergence, porabljen računski čas in natančnost izračunanih lastnih vrednosti za vsako metodo posebej. Analizirane metode so vse numerično stabilne, ekonomične in enostavno uporabne. Podprogrami v FORTRAN-u in primer glavnega programa so dodani za Mullerjevo in Laguerrovo metodo, ki sta se izkazali za bolj učinkoviti od Newtonove.

## Introduction

In this paper three numerical methods for the solution of the generalized algebraic eigenvalue problem

$$A(z)x = 0 \quad (1)$$

are described and the programmes of two of them are presented. Matrix  $A(z)$  is of the form

$$A(z) = A_0 + A_1 z + \dots + A_m z^m \quad (2)$$

where all  $A_i$  are real square matrices of order  $n$ .

The standard source of this problem is the solution of a system of  $n$  linear homogeneous ordinary differential equations of order  $m$ .

If the matrix  $A_m$  is non-singular then the problem (1) is equivalent to the classical eigenvalue problem

$$\begin{bmatrix} 0 & I & 0 & \dots & 0 \\ 0 & 0 & I & \dots & 0 \\ - & - & - & - & - \\ 0 & 0 & 0 & \dots & I \\ B_0 & B_1 & B_2 & \dots & B_{m-1} \end{bmatrix} \begin{bmatrix} y_0 \\ y_1 \\ - \\ y_{m-2} \\ y_{m-1} \end{bmatrix} = z \begin{bmatrix} y_0 \\ y_1 \\ - \\ y_{m-2} \\ y_{m-1} \end{bmatrix} \quad (3)$$

of order  $m \cdot n$  where  $B_i = -A_m^{-1} A_i$ . A similar reduction is possible if  $A_0$  is non-singular.

Although we have several efficient and stable numerical methods for the solution of the classical eigenvalue problem, for instance, QR algorithm [4], it is not economical to reduce the problem (1) to the standard form (3). Moreover, many practical problems like "flutter" problem in aerodynamics, are such that both  $A_0$  and  $A_m$  are singular and the reduction to the standard form causes further difficulties [3].

It is therefore considered as more advantageous to work directly with the matrix (2) in all cases. The most natural approach is to

determine the zeroes of the polynomial

$$f(z) = \det A(z) \quad (4)$$

using some of the many well known methods.

There are some efficient and stable algorithms for the evaluation of the determinant of the matrix, for instance, Gaussian elimination with pivoting [4]. It has been believed so far that the evaluation of the derivative of the determinant involves much more work than the evaluation of the determinant does, and therefore the interpolation methods which require only function values had been suggested for the solution of the algebraic equation

$$f(z) = 0 \quad (5)$$

see [3], [4]. The most popular of these methods is Muller's method [4] of successive quadratic interpolation. The prejudice against the methods which require the derivatives of the function is probably based on the fact that the derivative of the determinant is a sum of  $n$  determinants. However, there can be found in [1] a formula for the logarithmic derivative of the determinant which together with its derivative enable us to solve (5) by using Newton's or Laguerre's method efficiently.

In this paper algorithms for the solution of the generalized eigenvalue problem (1), based on Muller's, Newton's and Laguerre's methods are described and compared with regard to: (i) the number of operations (multiplications and divisions) involved in iteration step, (ii) the number of iterations, and (iii) the computer time used. The computations were performed on a number of typical examples.

#### Muller's method

Muller's method for the solution of the equation (5) as described in [4] is the following iterative process. From the three previous consecutive approximations to a root of (5):  $z_{r-2}$ ,  $z_{r-1}$ ,  $z_r$ , the parabola through the points

$$(z_{r-2}, f(z_{r-2})), (z_{r-1}, f(z_{r-1})), (z_r, f(z_r))$$

is formed and its zero which is the closest to  $z_r$  is accepted as the next approximation, say  $z_{r+1}$ , to a root. The asymptotic rate of convergence to a simple root is of order 1.84 and at each iteration step, only one new function value is required. It is necessary to use complex arithmetic throughout.

The algorithm is as follows.

For  $r=2, 3, \dots$  we calculate

$$h_r = z_r - z_{r-1}, \quad k_r = h_r/h_{r-1}, \quad d_r = 1 + k_r$$

$$g_r = f_{r-2}k_r^2 - f_{r-1}d_r^2 + f_r(k_r + d_r)$$

$$z_{r+1} = z_r - 2f_r h_r d_r / (g_r \pm (g_r^2 - 4f_r d_r k_r (f_{r-2}k_r - f_{r-1}d_r + f_r))^{1/2}) \quad (6)$$

$$f_{r+1} = f(z_{r+1})$$

The sign in the denominator of (6) is chosen so as to give the correction to  $z_r$  the smaller absolute magnitude.

The choice of the first three approximations is left to the user. There are no proofs of the global convergence of the iterative process.

In this case we have to calculate the value of

$$f = f(z) = \det A(z)$$

at each iteration step what can be done most efficiently by triangular decomposition of the matrix  $A(z)$  using partial pivoting as follows. For the given value of  $z$  we first calculate the elements of the matrix

$$A = A(z) \quad (7)$$

what requires  $m \cdot n^2$  operations. Then we decompose  $A$  into the form

$$PA = LU \quad (8)$$

where  $L$  is unit lower triangular matrix,  $U$  is upper triangular matrix, and  $P$  is the corresponding permutation matrix [4]. This step requires  $n^3/3 - n/3$  operations. Finally we compute

$$f = \pm \det U = \pm u_{11} u_{22} \dots u_{nn}$$

what requires  $n-1$  operations. Together we must perform

$$op_M = m \cdot n^2 + n^3/3 + 2n/3 - 1$$

operations for the calculation of  $f$ .

#### Newton's method

The well known Newton's method for the solution of (5) is the iterative process

$$z_{r+1} = z_r - f(z_r)/f'(z_r), \text{ for } r=1, 2, \dots$$

At each iteration step we evaluate the correction of the current approximation to a root as the reciprocal of the logarithmic derivative of the function. From [1], [2] we have the formula

$$f'(z)/f(z) = \text{tr}(A^{-1}(z)A'(z)) = \text{tr } X \quad (9)$$

where  $A'(z)$  denotes the derivative of the matrix  $A(z)$  with respect to  $z$ .

An efficient algorithm for calculating the expression (9) is as follows.

First we calculate the matrices

$$A = A(z), \quad B = A'(z) \quad (10)$$

which requires  $(2m-1)n^2$  operations. Then we decompose  $A$  as in (8) obtaining

$$PA = LU \quad (11)$$

which requires  $n^3/3 - n/3$  operations, compute the matrices  $Y$  and  $X$  which satisfy

$$LY = PB \quad (12)$$

$$UX = Y \quad (13)$$

As we need only the trace of  $X$ , we can use formulae

$$x_{ik} = (y_{ik} - \sum_{j=i+1}^n u_{ij}x_{jk})/u_{ii} \quad (14)$$

for  $i=n, n-1, \dots, 1$  and  $k=1, 2, \dots, i$ . These two steps require together  $(n^3/2 - n^2/2) + (n^3/6 + n^2/2 + n/3)$  operations. So we can calculate

$$f'(z)/f(z) = \sum_{i=1}^n x_{ii} \quad (15)$$

with

$$op_N = (2m-1)n^2 + n^3 \quad (16)$$

operations..

Although the asymptotic rate of convergence of the Newton's method to a simple root is 2 there seem to be little advantage in using Newton's method against Muller's. There appear to be almost three times as many operations per iteration step in Newton's method than in Muller's.

#### Laguerre's method

The quite famous Laguerre's method requires apart from the function value also the values of its first two derivatives. The asymptotic rate of convergence to a simple root is 3. Perhaps the most attractive feature from the users point of view is its stability in global convergence for any value of the first approximation.

The iterative process is defined by [4]

$$z_{r+1} = z_r - n/(S_1 \pm ((n-1)(nS_2 - S_1^2))^{1/2})$$

for  $r=1, 2, \dots$ , where

$$S_1 = f'(z_r)/f(z_r)$$

$$S_2 = (f''(z_r)^2 - f(z_r)f'''(z_r))/f(z_r)^2$$

The sign in the denominator is chosen so as to give the correction to  $z_r$  the smaller absolute magnitude. The form of  $S_1$  which is suitable for numerical evaluation is defined by (9) and (10) - (15), while the formula for  $S_2$  is derived from (9) by differentiation, see [2],

$$S_2 = \text{tr}((A^{-1}(z_r)A'(z_r))^2) - \text{tr}(A^{-1}(z_r)A''(z_r)) \quad (17)$$

The algorithm for evaluation of  $S_2$  demands in addition to the starting operations for  $S_1$ , i.e. (10) and (11), the following computations:

We must first calculate the matrix

$$C = A''(z) \quad (18)$$

with additional  $(m-2)n^2$  operations. Then we calculate the trace of the matrix  $W = A^{-1}C$  by steps

$$LZ = PC \quad (19)$$

$$UW = Z \quad (20)$$

$$\text{tr } W = \sum_{i=1}^n w_{ii} \quad (21)$$

which require  $(n^3/2 - n^2/2) + (n^3/6 + n^2/2 + n/3) + (n-1)$  operations.

The first term in (17) is found by calculating first the full matrix  $X$  by (12) and (13) with  $n^3$  operations and then the trace of  $X^2$  as

$$\text{tr}(X^2) = \sum_{i=1}^n \sum_{k=1}^n x_{ik}x_{ki} \quad (22)$$

which requires additional  $n^2$  operations.

At each iteration step we must perform

$$op_L = (3m-3)n^2 + 2n^3 + n^2$$

operations for the calculation of  $S_1$  and  $S_2$ . This number is almost twice as large as (16) for the Newton's method.

#### Numerical tests

The programmes in FORTRAN programming language were written for all three algorithms and tested on CDC CYBER 72 computer. Single precision complex arithmetic was used with  $t=48$ , where  $t$  is the number of significant digits in the mantissa of a binary floating-point number. The iterative process was stopped after either the computed correction to the computed approximation to the eigenvalue was less than or equal to  $\epsilon$  or if  $|u_{ii}| \leq \epsilon$ , for  $1 \leq i \leq n$ , occurred in (8) or (11). The value for  $\epsilon$  was chosen as  $10 \cdot \|A(z)\|_{\infty} \cdot 2^{-t}$ .

The implicit deflation of  $f(z)$  was applied in the programmes in order to prevent the iterative process to converge to one of the previously computed eigenvalues, say  $x_1, x_2, \dots, x_m$ . This introduces the following modifications of algorithms:

$$f(z) \text{ into } f(z) / \prod_{i=1}^m (z-x_i) \quad (23)$$

$$f'(z)/f(z) \text{ into } f'(z)/f(z) - \sum_{i=1}^m (z-x_i)^{-1} \quad (24)$$

$$S_2 \text{ into } S_2 - \sum_{i=1}^m (z-x_i)^{-2} \quad (25)$$

Some typical examples of (1) were computed with multiple eigenvalues either of finite value or in the infinity. The values of the parameters  $m$  and  $n$  of  $A(z)$  were on intervals  $1 \leq m \leq 2$  and  $3 \leq n \leq 12$  respectively. No breakdown of the iterative process occurred regardless the algorithm used. All three methods proved to be stable and accurate. It was found that the previous computed eigenvalue and the values round it were the most suitable starting values for computing the next eigenvalue.

In the table

| i   | Muller |     | Newton |      | Laguerre |     |
|-----|--------|-----|--------|------|----------|-----|
|     | No     | T   | No     | T    | No       | T   |
| 1   | 55     | .57 | 41     | .58  | 17       | .45 |
| 2   | 0      |     | 0      |      | 0        |     |
| 3   | 11     | .13 | 19     | .28  | 11       | .30 |
| 4   | 0      |     | 0      |      | 0        |     |
| 5   | 3      | .04 | 4      | .06  | 2        | .06 |
| 6   | 0      |     | 0      |      | 0        |     |
| 7   | 16     | .17 | 41     | .61  | 3        | .08 |
| 8   | 1      | .02 | 2      | .04  | 2        | .06 |
| ALL | 86     | .93 | 107    | 1.57 | 35       | .95 |

the number of iteration steps (No) and the computer time used (T in sec) are shown for each of the computed eigenvalue  $z_i, i=1, \dots, 8$ , of the system

$$(A_0 + zA_1 + z^2A_2)x = 0$$

where  $A_2 = I$ , and

$$A_1 = \begin{bmatrix} 0 & -3 & 0 & -1 \\ 2 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 \\ 0 & 0 & 2 & 0 \end{bmatrix}, \quad A_0 = \begin{bmatrix} -1 & 0 & 0 & 0 \\ 0 & -2 & 0 & -1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{bmatrix}.$$

The system has a triple eigenvalue at  $\pm i$  and a double eigenvalue at 0 [5]. Somehow similar results were obtained with other examples.

Three more examples are

$$1) A(z) = A_0 + A_1z + A_2z^2$$

$$A_0 = \begin{bmatrix} 1 & -1 & 1 \\ -15 & 0 & 0 \\ 1 & 0 & 1 \end{bmatrix}, \quad A_1 = \begin{bmatrix} -2 & 1 & -1 \\ 3 & 0 & 1 \\ 1 & 0.5 & 0 \end{bmatrix},$$

$$A_2 = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 0.25 & 0 \\ -1 & 0 & 1 \end{bmatrix}$$

$$z_i:$$

1.000000000000  
2.08866333896  
-0.256555796702 + 0.896010203022 i  
-2.91609433059  
11.3405425851

$$2) A(z) = A_0 + A_1z$$

$$A_0 = \begin{bmatrix} -1 & -3 & -3 & -3 & -3 & -3 \\ -3 & -4 & -3.1 & -3.1 & -3.1 & -3.1 \\ -3 & -3.1 & 2.8 & 3.8 & 3.8 & 3.8 \\ -3 & -3.1 & 3.8 & 9.8 & 10.7 & 10.7 \\ -3 & -3.1 & 3.8 & 10.7 & 12.6 & 14.6 \\ -3 & -3.1 & 3.8 & 10.7 & 14.6 & 15.6 \end{bmatrix}$$

$$A_1 = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & -1 & -1 & -1 & -1 \\ 1 & 0 & -1 & -2 & -2 & -2 \\ 1 & 0 & -1 & -2 & -3 & -3 \\ 1 & 0 & -1 & -2 & -3 & -2 \end{bmatrix}$$

$$z_i:$$

0.908770404173 + 1.93967680102 i  
0.931536974557 + 1.97197662562 i  
4.18245919165  
6.13692605089

$$3) A(z) = A_0 + A_1z + A_2z^2$$

$$A_0 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ -1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & -1 & 3 & 0 & 0 & 0 & 0 \\ 0 & 0 & -4 & 7 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1.5 & 1 & 0 & 0 \\ 0 & 0 & -8 & 0 & 0 & 10 & 0 \\ 0 & 0 & 0 & 0 & 0 & -0.4 & 1 \end{bmatrix}$$

$$A_1 = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & -1 & 2 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 6 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & -1 & 0 & 0 & 9 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$A_2 = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 5 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$z_1:$

-1.09579878411  
 -1.08865049880 ± 1.04888469390 i  
 -0.541227886923 ± 1.51715942168 i

$$z_6 - z_{14} \approx K \cdot 10^7, \quad |K| \approx 1$$

The last nine eigenvalues lie in the infinity.

Muller's and Laguerre's methods were also tested on DEC 1091 computer with  $t=27$ . In one case when some of the elements of  $A(z)$  differed up to  $10^{13}$  in their absolute values Muller's method failed to compute correctly some of the eigenvalues while all the eigenvalues obtained by Laguerre's method were correct.

#### References

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#### Programmes

##### 1. Comments

In this chapter two subroutines EIGENM and EIGENL for computing the eigenvalues of a general matrix (2) are presented. An example of a possible main (calling) programme for both subroutines is also given at the end of the chapter. In EIGENM the eigenvalues are computed by Muller's method while in EIGENL by Laguerre's method. If  $A(z)$  in (2) satisfies  $m \leq 3$  then the subroutines can be used such as they are presented. For  $m > 3$  some changes in the subroutines are necessary as explained at the end of the programmes.

The subroutines EIGENM and EIGENL are completely selfcontained. (EIGENM is composed of four subroutines EIGENM, F, ALAMDA, LU and EIGENL is composed of ten subroutines EIGENL, ALAMDA, LU, DERIV1, DERIV2, LOWER, UPPER, UPPER2, TRACE, TRXSQ) and communication to them is solely through their argument lists and COMMON statements. The entrance to the subroutine EIGENM is achieved by

```
COMMON /MATRIX/ AO(ND,ND), A1(ND,ND), A2(ND,
1ND)
```

```
CALL EIGENM (M, N, ND, T, INDIC, LAM, PERMUT,
1A, L, U)
```

The entrance to the subroutine EIGENL is achieved by

```
COMMON /MATRIX/ AO(ND, ND), A1(ND,ND), A2(ND,
1ND), A3(ND,ND)
```

```
CALL EIGENL (M, N, ND, T, INDIC, LAM, PERMUT,
1A, D, L, U)
```

The meaning of the parameters is described in the following lines.

The elements of matrices  $A_0, A_1, A_2, \dots, A_m$  in (2) are to be stored in the first N rows and columns of the two dimensional arrays  $A_0, A_1, A_2, A_3, \dots$ .

M, where  $M \geq 1$ , is the degree of z-matrix  $A(z)$  in (2).

N, where  $N \geq 1$ , is the order of matrices  $A_0, A_1, A_2, \dots, A_m$ .

ND, where  $ND \geq N$ , defines the first dimension of the two dimensional arrays  $A_0, A_1, A_2, \dots, A, D, L, U$  and the dimension of the one dimensional array PERMUT. Dimension of the one dimensional arrays INDIC and LAM must be equal to or greater than  $ND \cdot M$ .

T is the number of binary digits in the

mantissa of a single precision floating-point number. T is of integer type.

The array INDIC indicates the success of the subroutine as follows

| value of INDIC(I) | eigenvalue (I)                      |
|-------------------|-------------------------------------|
| 0                 | not found after 100 iteration steps |
| 1                 | found                               |

LAM is one dimensional complex array. The computed eigenvalues will be found in the first N\*M places.

The arrays PERMUT, A, D, L, U is the working storage.

The main (calling) programme should contain the following declaration statements

```

INTEGER, M,N,T,INDIC,PERMUT
REAL A0,A1,A2,A3
COMPLEX LAM,A,L,U
COMMON /MATRIX/ A0(ND,ND), A1(ND,ND),
1A2(ND,ND),A3(ND,ND)
DIMENSION INDIC (ND*M),LAM(ND*M),PERMUT
1(ND),A(ND,ND+1),L(ND,ND+1),U(ND,ND+1)
EQUIVALENCE (A(1),L(1),U(1))

```

and the following two additional declaration statements for the subroutine EIGENL

```

COMPLEX D
DIMENSION D(ND,ND)

```

Note that the dimensions of the arrays in DIMENSION and COMMON statements of the main programme and in COMMON statements within subroutines must be constants of integer type. See the example in 5 where ND was replaced by 12, ND+1 by 13 and ND\*M by 36.

## 2. Subroutine EIGENM - Muller's method

The subroutine is composed of four subroutines EIGENM, F, ALAMDA and LU.

```

SUBROUTINE EIGENM(M,N,ND,T,INDIC,LAM,
1PERMUT,A,L,U)
INTEGER M,N,ND,T,INDIC,PERMUT,I,IND,
1ITERST, NM, R
REAL EPS, EPS100, IM, RE
REAL A0,A1,A2,A3
COMPLEX LAM,A,L,U,DENOM1,DENOM2,DR,FR,
1FRM1,FRM2,GR,HR,HRM1,KR,ZR,ZRM1,ZRM2
COMMON /MATRIX/ A0(ND,ND),A1(ND,ND),

```

```

1A2(ND,ND),A3(ND,ND)
DIMENSION INDIC(1),LAM(1),PERMUT(1),
1A(ND,1),L(ND,1),U(ND,1)
NM = N*M
ITERST = 102
DO 1 I=1,NM
    INDIC(I) = 0

```

1 CONTINUE

I = 1

C THE MULLER ITERATIVE PROCESS.

C DETERMINATION OF THREE STARTING

C CONSECUTIVE APPROXIMATIONS TO THE

C COMPUTED EIGENVALUE LAM(I),FOR

C I=1,2,...,NM.

2 ZR = (1.0,0.0)\*FL0AT(I)

3 ZRM1 = 0.5\*ZR

ZRM2 = 0.1\*ZR

C THE INITIAL COMPUTATIONS OF THE NEXT

C APPROXIMATION Z(R+1) TO LAM(I),FOR R=2.

R = 2

CALL F(ZRM2,FRM2,EPS,I,IND,M,N,ND,T,LAM,

1PERMUT,A,L,U)

IF(IND.EQ.1)G0 TO 5

CALL F(ZRM1,FRM1,EPS,I,IND,M,N,ND,T,LAM,

1PERMUT,A,L,U)

IF(IND.EQ.1)G0 TO 6

CALL F(ZR,FR,EPS,I,IND,M,N,ND,T,LAM,

1PERMUT,A,L,U)

IF(IND.EQ.1)G0 TO 7

C

HRM1 = ZRM1-ZRM2

HR = ZR -ZRM1

KR = HR/HRM1

C

THE MAIN LOOP FOR THE COMPUTATION OF

C

THE NEXT APPROXIMATION Z(R+1) TO LAM(I),

C

FOR R=2,3,...,ITERST.

4 DR=1.0+KR

GR=FRM2\*KR\*\*2 - FRM1\*DR\*\*2 + FR\*(KR+DR)

DENOM2 = CSQRT(GR\*\*2-4.\*FR\*DR\*KR)

1 (FRM2\*KR-FRM1\*DR+FR)

DENOM1 = GR+DENOM2

DENOM2 = GR-DENOM2

IF(CABS(DENOM2).GT.CABS(DENOM1))DENOM1

1 = DENOM2

KR = -2.0\*FR\*DR/DENOM1

HR = HR\*KR

GR = ZR

ZR = ZR+HR

FRM2 = FRM1

FRM1 = FR

CALL F(ZR,FR,EPS,I,IND,M,N,ND,T,LAM,

1 PERMUT,A,L,U)

C

IF(IND.EQ.1)G0 TO 7

IF(CABS(ZR-GR).LE.10.\*EPS)G0 TO 7

```

      R = R+1
      IF(R.LT.ITERST)G0 T0 4
      LAM(I) = ZR
      G0 T0 9
5  ZR = ZRM2
      G0 T0 7
6  ZR = ZRM1
7  LAM(I) = ZR
      INDIC(I) = 1
      I = I+1
C   COMPUTE THE COMPLEX CONJUGATE VALUE OF
C   LAM(I).
      RE = REAL(ZR)
      IM = AIMAG(ZR)
      EPS100 = 100.0*EPS
      IF(ABS(IM).LE.EPS100)G0 T0 8
      LAM(I) = CONJG(ZR)
      INDIC(I) = 1
      I = I+1
C
8  IF(NM.LT.I)G0 T0 9
C   DETERMINATION OF THE STARTING APPROXIM
C   ATION TO THE COMPUTED EIGENVALUE LAM
C   (I+1).
      ZR = 0.9*ZR
      RE = ABS(RE)+ABS(IM)
      IF(RE.LE.EPS100.0R.1./RE.LE.EPS100)G0
1T0 2
      G0 T0 3
C
9  CONTINUE
      RETURN
      END

      SUBROUTINE F(Z,FUNC,EPS,I,IND,M,N,ND,T,
1LAM,PERMUT,A,L,U)
      INTEGER I,IND,M,N,ND,T,PERMUT,IM1,J,K
      REAL EPS,FK
      COMPLEX Z,FUNC,LAM,A,L,U,C
      DIMENSION LAM(1),PERMUT(1),A(ND,1),
1L(ND,1),U(ND,1)
C
C   COMPUTATION OF FUNC = F(Z) = DET A(Z) =
C   DET A. MATRIX A IS DECOMPOSED INTO THE
C   FORM
C
C       PA = LU
C
C   WHERE L IS UNIT LOWER TRIANGULAR MATRIX,
C   U IS UPPER TRIANGULAR MATRIX AND P IS
C   THE CORRESPONDING PERMUTATION MATRIX
C   (SEE WILKINSON).
C
C   CALCULATION OF A(Z).
      CALL ALAMDA (Z,M,N,ND,A)
C   CALCULATION OF THE SMALL POSITIVE
C   NUMBER EPS.
      C = (0.,0.)

```

```

      D0 1 J=1,N
      D0 1 K=1,N
      C = C + A(J,K)**2
1  CONTINUE
      EPS = CABS(CSQRT(C))*2.0**(-T)
C   COMPUTATION OF THE MATRICES L AND U,
C   WHERE PA = L*U.
      CALL LU(EPS, IND,N,ND,PERMUT,A,L,U)
      IF(IND.EQ.1)G0 T0 4
C
      J = -IND
      K = (-1)**J
      FK = FLOAT(K)
      FUNC = CMPLX(FK,0.0)
      D0 2 K=1,N
      FUNC = FUNC*U(K,K)
2  CONTINUE
C   MODIFICATION OF FUNC = F(Z) INTO F(Z)/
C   ((Z-X(1))(Z-X(2))...(Z-X(I-1))) ACCORDING
C   TO EQUATION (23) OF THIS PAPER.
      IM1 = I-1
      IF(IM1.EQ.0)G0 T0 4
      D0 3 J=1,IM1
      FUNC = FUNC/(Z-LAM(J))
3  CONTINUE
4  CONTINUE
      RETURN
      END

      SUBROUTINE ALAMDA(Z,M,N,ND,A)
      INTEGER M,N,ND,I,J,K
      REAL A0,A1,A2,A3
      COMPLEX Z,A,AIJ
      COMPLEX Z2,Z3
      COMMON /MATRIX/ A0(ND,ND),A1(ND,ND),
1A2(ND,ND),A3(ND,ND)
      DIMENSION A(ND,1)
C
C   THIS SUBROUTINE CALCULATES THE ELEMENTS
C   OF MATRIX A(Z), WHICH ARE FUNCTIONS OF Z,
C   FOR A GIVEN VALUE OF Z.
      Z2 = Z**2
      Z3 = Z**3
      CONTINUE
      D0 44 I=1,N
      D0 44 J=1,N
      AIJ = A0(I,J)
      D0 33 K=1,M
      G0 T0 (1,2,3),K
1      AIJ = AIJ+A1(I,J)*Z
      G0 T0 33
2      AIJ = AIJ+A2(I,J)*Z2
      G0 T0 33
3      AIJ = AIJ+A3(I,J)*Z3
33  CONTINUE
      A(I,J) = AIJ

```

```

44 CONTINUE
RETURN
END

```

```

SUBROUTINE LU (EPS,IND,N,ND,PERMUT,A,L,
1U)
INTEGER IND,N,ND,PERMUT,I,K,NP1,R,RC,
1RM1,RP1,T
REAL EPS,SRC,STABS
COMPLEX A,L,U,ST,URR
DIMENSION PERMUT(1),A(ND,1),L(ND,1),
1U(ND,1)

```

```

C
C THIS SUBROUTINE PERFORMS THE
C DECOMPOSITION OF MATRIX (PA) INTO L*U,
C WHERE L IS UNIT LOWER TRIANGULAR MATRIX,
C U IS UPPER TRIANGULAR MATRIX AND P IS
C THE CORRESPONDING PERMUTATION MATRIX.
C PARTIAL PIVOTING IS INTRODUCED (SEE J.H.
C WILKINSON, PAGE 225).

```

```
IND = 0
```

```
DO 1 I=1,N
```

```
PERMUT(I) = I
```

```
1 CONTINUE
```

```
C
```

```
NP1 = N+1
```

```
DO 11 R=1,N
```

```
RM1 = R-1
```

```
RP1 = R+1
```

```
RC = R
```

```
SRC = 0
```

```
DO 4 T=R,N
```

```
ST = A(T,R)
```

```
IF(RM1.EQ.0)GO TO 3
```

```
C
```

```
DO 2 K=1,RM1
```

```
ST = ST-L(T,K)*U(K,R)
```

```
2 CONTINUE
```

```
3 A(T,NP1) = ST
```

```
STABS = CABS(ST)
```

```
IF(SRC.GE.STABS)GO TO 4
```

```
SRC = STABS
```

```
RC = T
```

```
4 CONTINUE
```

```
IF(SRC.GT.EPS)GO TO 5
```

```
IND = 1
```

```
RETURN
```

```
C
```

```
5 IF(R.EQ.RC)GO TO 7
```

```
T = PERMUT(R)
```

```
PERMUT(R) = PERMUT(RC)
```

```
PERMUT(RC) = T
```

```
C
```

```
THE FOLLOWING STATEMENT(IND=IND-1) MAY
BE TAKEN OUT IN EIGENL.
```

```
IND = IND-1
```

```
DO 6 K=1,NP1
```

```
ST = A(R,K)
```

```
A(R,K) = A(RC,K)
```

```
A(RC,K) = ST
```

```
6 CONTINUE
```

```
7 URR = A(R,NP1)
```

```
U(R,R) = URR
```

```
IF(R.EQ.N)GO TO 11
```

```
DO 10 T=RP1,N
```

```
L(T,R) = A(T,NP1)/URR
```

```
ST = A(R,T)
```

```
IF(RM1.EQ.0)GO TO 9
```

```
C
```

```
DO 8 K=1,RM1
```

```
ST = ST - L(R,K)*U(K,T)
```

```
8 CONTINUE
```

```
9 U(R,T) = ST
```

```
10 CONTINUE
```

```
11 CONTINUE
```

```
RETURN
```

```
END
```

### 3. Subroutine EIGENL - Laguerre's method

The subroutine is composed of ten subroutines  
EIGENL,ALAMDA (see 2.), LU(see 2.), DERIV1,  
DERIV2, LOWER, UPPER, UPPER2, TRACE and TRXSQ.

```

SUBROUTINE EIGENL(M,N,ND,T,INDIC,LAM,
1PERMUT,A,D,L,U)
INTEGER M,N,ND,T,INDIC,PERMUT,I,IM1,IND,
1ITERST,J,K,NM,NM1,R
REAL EPS,EPS100,IM,RE
REAL A0,A1,A2,A3
COMPLEX LAM,A,D,L,U,CEPS,DENOM1,DENOM2,
1S1,S2,TRACEX,TRACEW,TRACX2,ZR,ZRP1
COMMON /MATRIX/ A0(ND,ND),A1(ND,ND)
1A2(ND,ND),A3(ND,ND)
DIMENSION INDIC(1),LAM(1),PERMUT(1),
1A(ND,1),D(ND,1),L(ND,1),U(ND,1)

```

```
C
```

```
NM = N*M
```

```
ITERST = 101
```

```
DO 1 I=1,NM
```

```
INDIC(I)=0
```

```
1 CONTINUE
```

```
I=1
```

```
C
```

```
THE LAGUERRE ITERATIVE PROCESS.
```

```
C
```

```
DETERMINATION OF THE STARTING
```

```
C
```

```
APPROXIMATION TO THE FIRST COMPUTED
```

```
C
```

```
EIGENVALUE LAM(1).
```

```
ZR = (1.,1.)
```

```
2 NM1 = NM-I
```

```
C
```

```
COMPUTATION OF THE NEXT EIGENVALUE LAM(I).
```

```
R = 1
```

```
C
```

```
THE MAIN LOOP FOR THE COMPUTATION OF THE
```



```

C   NEXT APPROXIMATION Z(R+1) TO LAM(I).
C   CALCULATION OF THE EQUATIONS (10) - (22)
C   OF THIS PAPER.
C   CALCULATION OF A(Z).
3  CALL ALAMDA(ZR,M,N,ND,A)
C   CALCULATION OF THE SMALL POSITIVE NUMBER
C   EPS.
    CEPS = 0
    DO 4 J=1,N
      DO 4 K=1,N
        CEPS = CEPS+A(J,K)**2
4  CONTINUE
    EPS = CABS(CSQRT(CEPS))*2.0**(-T)
    EPS100 = 100.*EPS
C   COMPUTATION OF THE MATRICES L AND U,
C   WHERE PA = L*U.
    CALL LU(EPS,IND,N,ND,PERMUT,A,L,U)
    IF(IND.EQ.1)GO TO 8
C   COMPUTATION OF THE TRACE OF MATRIX X
C   AND THE TRACE OF X**2.
    CALL DERIV1(ZR,M,N,ND,PERMUT,D)
    CALL LOWER(N,ND,D,L)
    CALL UPPER(N,ND,D,U)
    CALL TRACE(N,ND,D,TRACEX)
    CALL TRXSQ(N,ND,D,TRACX2)
C   COMPUTATION OF THE TRACE OF MATRIX W.
    TRACEW = 0.0
    IF(M.LT.2)GO TO 5
    CALL DERIV2(ZR,M,N,ND,PERMUT,D)
    CALL LOWER(N,ND,D,L)
    CALL UPPER2(N,ND,D,U)
    CALL TRACE(N,ND,D,TRACEW)
C
5  S1 = TRACEX
   S2 = TRACX2-TRACEW
C   MODIFICATIONS OF S1 AND S2 ACCORDING
C   TO THE EQUATIONS (24) AND (25) OF THIS
C   PAPER.
    IM1 = I-1
    IF(IM1.EQ.0)GO TO 7
    DO 6 J=1,IM1
      CEPS = 1.0/(ZR-LAM(J))
      S1 = S1-CEPS
      S2 = S2-CEPS**2
6  CONTINUE
C
7  CEPS = CSQRT(MM1*((NM1+1)*S2-S1**2))
   DENOM1 = S1+CEPS
   DENOM2 = S1-CEPS
   IF(CABS(DENOM1).GT.CABS(DENOM2))DENOM2=
1DENOM1
   ZRP1 = ZR-(NM1+1)/DENOM2
C
   IF(CABS(ZRP1-ZR).LE.10.*LPS)GO TO 9
   ZR = ZRP1
   R = R+1

```

```

   IF(R.LT.ITERST)GO TO 3
   LAM(I) = ZR
   GO TO 11
C
8  ZRP1 = ZR
9  LAM(I) = ZRP1
   INDIC(I) = 1
C   DETERMINATION OF THE STARTING
C   APPROXIMATION TO THE NEXT COMPUTED
C   EIGENVALUE LAM(I+1) AND COMPUTATION OF
C   THE COMPLEX CONJUGATE VALUE OF LAM(I).
   ZR = 0.9*ZRP1
   I = I+1
   RE = REAL(ZRP1)
   IM = AIMAG(ZRP1)
   IF(ABS(IM).LE.1.0E-3*ABS(RE))ZR=0.5*
1CMPLX(RE,RE)
   RE = ABS(RE)+ABS(IM)
   IF(RE.LE.EPS100.0R.1./RE.LE.EPS100)
1ZR=(1.,1.)*FL0AT(I)
   IF(ABS(IM).LE.EPS100)GO TO 10
   LAM(I) = CONJG(ZRP1)
   INDIC(I) = 1
   I = I+1
C
10 IF(I.LE.NM)GO TO 2
11 CONTINUE
   RETURN
   END

SUBROUTINE DERIV1(Z,M,N,ND,PERMUT,A)
  INTEGER M,N,ND,PERMUT,I,II,J,K,MM1
  REAL A0,A1,A2,A3
  COMPLEX Z,A,AIJ
  COMPLEX Z2,Z3
  COMMON /MATRIX/ A0(ND,ND),A1(ND,ND)
  1A2(ND,ND),A3(ND,ND)
  DIMENSION A(ND,1),PERMUT(1)
C
C   THIS SUBROUTINE CALCULATES THE ELEMENTS
C   OF THE FIRST DERIVATIVE OF MATRIX A(Z),
C   FOR A GIVEN VALUE OF Z.
   Z2 = 2.*Z
   Z3 = 3.*Z**2
  CONTINUE
  MM1 = M-1
  DO 55 II=1,N
    I = PERMUT(II)
    DO 55 J=1,N
      AIJ = A1(I,J)
      IF(M.LT.2)GO TO 44
      DO 33 K=1,MM1
        GO TO (2,3),K
      2  AIJ = AIJ+A2(I,J)*Z2
        GO TO 33
      3  AIJ = AIJ+A3(I,J)*Z3

```

```

33      CONTINUE
44      A(II,J)=AIJ
55 CONTINUE
      RETURN
      END

      SUBROUTINE DERIV2(Z,M,N,ND,PERMUT,A)
      INTEGER M,N,ND,PERMUT,I,II,J,K,MM2
      REAL A0,A1,A2,A3
      COMPLEX Z,A,AIJ
      COMPLEX Z2,Z3
      COMMON /MATRIX/ A0(ND,ND),A1(ND,ND),
1A2(ND,ND),A3(ND,ND)
      DIMENSION A(ND,1),PERMUT(1)

C
C      THIS SUBROUTINE CALCULATES THE ELEMENTS
C      OF THE SECOND DERIVATIVE OF MATRIX A(Z),
C      FOR A GIVEN VALUE OF Z.
      Z3 = 6.*Z
      CONTINUE
      MM2 = M-2
      DO 55 II=1,N
          I = PERMUT(II)
          DO 55 J=1,N
              AIJ = 2.*A2(I,J)
              IF(M.LT.3)GO TO 44
              DO 33 K=1,MM2
                  GO TO (3,3),K
              3      AIJ = AIJ+A3(I,J)*Z3
33      CONTINUE
44      A(II,J) = AIJ
55 CONTINUE
      RETURN
      END

      SUBROUTINE LOWER(N,ND,A,L)
      INTEGER N,ND,I,IM1,J,K
      COMPLEX A,L,AIJ
      DIMENSION A(ND,1),L(ND,1)

C
C      THIS SUBROUTINE CALCULATES THE ELEMENTS
C      OF MATRIX M WHICH SATISFIES THE RELATION
C      L*M=A, WHERE L IS UNIT LOWER TRIANGULAR
C      MATRIX. M IS STORED INTO ARRAY A.
      DO 2 I=2,N
          IM1 = I-1
          DO 2 J=1,N
              AIJ = (0.,0.)
              DO 1 K=1,IM1
                  AIJ = AIJ+L(I,K)*A(K,J)
1      CONTINUE
              A(I,J)=A(I,J)-AIJ
2 CONTINUE
      RETURN
      END

```

```

      SUBROUTINE UPPER(N,ND,A,U)
      INTEGER N,ND,I,II,IP1,J,K
      COMPLEX A,U,AIJ
      DIMENSION A(ND,1),U(ND,1)

C
C      THIS SUBROUTINE CALCULATES THE ELEMENTS
C      OF MATRIX M WHICH SATISFIES THE RELATION
C      U*M=A, WHERE U IS UPPER TRIANGULAR
C      MATRIX. M IS STORED INTO ARRAY A.
      DO 3 II=1,N
          I = N+1-II
          IP1 = I+1
          DO 3 J=1,N
              AIJ = (0.,0.)
              IF(I.EQ.N)GO TO 2
              DO 1 K=IP1,N
                  AIJ = AIJ+U(I,K)*A(K,J)
1      CONTINUE
              A(I,J) = (A(I,J)-AIJ)/U(I,I)
3 CONTINUE
      RETURN
      END

      SUBROUTINE TRACE(N,ND,A,TR)
      INTEGER N,ND,I
      COMPLEX A,TR
      DIMENSION A(ND,1)

C
C      THIS SUBROUTINE CALCULATES THE TRACE
C      OF MATRIX A.
      TR = (0.,0.)
      DO 1 I=1,N
          TR = TR+A(I,I)
1 CONTINUE
      RETURN
      END

      SUBROUTINE TRXSQ(N,ND,A,TR)
      INTEGER N,ND,I,K
      COMPLEX A,TR
      DIMENSION A(ND,1)

C
C      THIS SUBROUTINE CALCULATES THE TRACE OF
C      A**2.
      TR = (0.,0.)
      DO 1 I=1,N
          DO 1 K=1,N
              TR = TR+A(I,K)*A(K,I)
1 CONTINUE
      RETURN
      END

      SUBROUTINE UPPER2 (N,ND,A,U)
      INTEGER N,ND,I,II,IP1,J,K
      COMPLEX A,U,AIJ
      DIMENSION A(ND,1),U(ND,1)

```

C THIS SUBROUTINE CALCULATES THE  
C ELEMENTS OF THE LOWER TRIANGULAR OF  
C MATRIX M. M SATISFIES THE RELATION  $U^*M=A$ ,  
C WHERE U IS UPPER TRIANGULAR MATRIX. M  
C IS STORED INTO ARRAY A.

```

D0 3 II=1,N
  I = N+1-II
  IP1 = I+1
  D0 3 J=1,I
    AIJ = (0.,0.)
    IF(I.EQ.N)G0 T0 2
    D0 1 K=IP1,N
      AIJ = AIJ+U(I,K)*A(K,J)
1    CONTINUE
2    A(I,J)=(A(I,J)-AIJ)/U(I,I)
3  CONTINUE
  RETURN
  END

```

#### 4. Modifications

For matrix  $A(z)$  satisfying (2), where  $m \geq 3$ , some of the statements must be replaced with other statements as shown below for the case  $m \leq 5$ :

##### 1) Statements

```

REAL A0,A1,A2,A3
COMMON /MATRIX/ A0(ND,ND),A1(ND,ND),
1A2(ND,ND),A3(ND,ND)

```

in the main programme, EIGENM, ALAMDA, EIGENL, DERIV1 and DERIV2 by

```

REAL A0,A1,A2,A3,A4,A5
COMMON /MATRIX/ A0(ND,ND),A1(ND,ND),
1A2(ND,ND),A3(ND,ND),A4(ND,ND),A5(ND,ND)

```

2) Statement COMPLEX Z2,Z3 in ALAMDA,DERIV1 and DERIV2 by

```
COMPLEX Z2,Z3,Z4,Z5
```

3) Statement CONTINUE

(i) in ALAMDA by

```

Z4 = Z**4
Z5 = Z**5

```

(ii) in DERIV1 by

```

Z4 = 4.*Z**3
Z5 = 5.*Z**4

```

(iii) in DERIV2 by

```

Z4 = 4.*3.*Z**2
Z5 = 5.*4.*Z**3

```

4) Statement  $G0 T0(1,2,3),K$  in ALAMDA by  
 $G0 T0(1,2,3,4,5),K$

5) Statement  $G0 T0(2,3),K$  in DERIV1 by  
 $G0 T0(2,3,4,5),K$

6) Statement  $G0 T0(3,3),K$  in DERIV2 by  
 $G0 T0(3,4,5),K$

7) Statement 33 CONTINUE in ALAMDA,DERIV1 and DERIV2 by

```

G0 T0 33
4 AIJ = AIJ+A4(I,J)*Z4
G0 T0 33
5 AIJ = AIJ+A5(I,J)*Z5
33 CONTINUE

```

#### 5. Main - calling programme

In the following lines an example of a possible calling programme for subroutines EIGENM and EIGENL is presented. It finds the eigenvalues of  $A(z)$  in (2) with  $m \leq 3$  and  $n \leq 12$ . The programme stops after reading value zero for n.

```

INTEGER INDIC,PERMUT,M,N,T,I,J,K,MN,MP1
REAL A0,A1,A2,A3
COMPLEX LAM,A,L,U
COMMON /MATRIX/ A0(12,12),A1(12,12),
1A2(12,12),A3(12,12)
EQUIVALENCE (A(1),L(1),U(1))
DIMENSION INDIC(36),LAM(36),PERMUT(12),
1A(12,13),L(12,13),U(12,13)
COMPLEX D
DIMENSION D(12,12)

```

C

```
1 READ 2,N,M,T
```

```
2 FORMAT(3I5)
```

```
3 FORMAT(8F10.0)
```

```
IF(N.EQ.0)G0 T0 12
```

C THE PROGRAMME STOPS AFTER READING ZERO

C FOR N.

```
MN = M*N
```

```
MP1 = M+1
```

```
D0 8 I=1,MP1
```

```
G0 T0 (4,5,6,7),I
```

```
4 READ 3,((A0(K,J),J=1,N),K=1,N)
```

```
G0 T0 8
```

```
5 READ 3,((A1(K,J),J=1,N),K=1,N)
```

```
G0 T0 8
```

```
6 READ 3,((A2(K,J),J=1,N),K=1,N)
```

```
G0 T0 8
```

```
7 READ 3,((A3(K,J),J=1,N),K=1,N)
```

```
8 CONTINUE
```

```

C
PRINT 9,((A0(K,J),J=1,N),K=1,N)
9 FORMAT (1H ,6E16.7)
PRINT 9,((A1(K,J),J=1,N),K=1,N)
IF(M.GT.1)PRINT 9,((A2(K,J),J=1,N),
1K=1,N)
IF(M.GT.2)PRINT 9,((A3(K,J),J=1,N),
1K=1,N)
C
CALL EIGENM(M,N,12,T,INDIC,LAM,PERMUT,
1A,L,U)
C
PRINT 10,(INDIC(I),LAM(I),I=1,MN)
10 FORMAT(1H0,10X,6H INDIC,10X,9HREAL PART,
112X,10H IMAG.PART,///(1H ,12X,I2,5X,
1E19.12,4X,E19.12))
D0 11 I=1,MN
LAM(I) = 0.0
11 CONTINUE
C
CALL EIGENL(M,N,12,T,INDIC,LAM,PERMUT,
1A,D,L,U)
PRINT10,(INDIC(I),LAM(I),I=1,MN)
G0 T0 1
C
12 CONTINUE
END

```

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